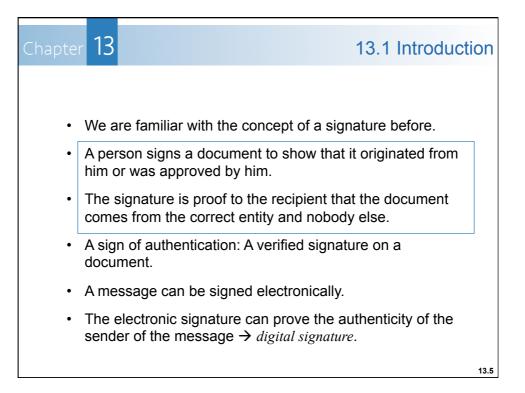
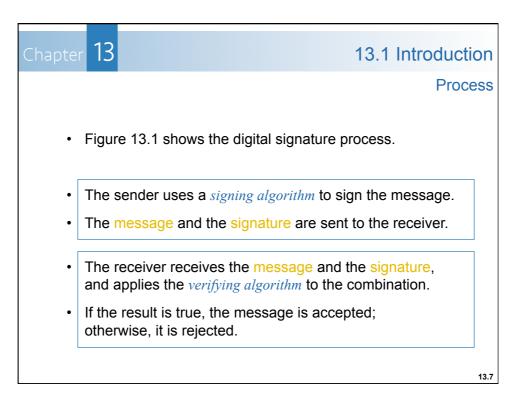
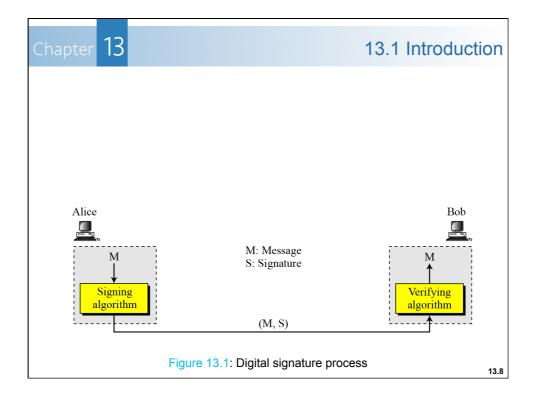


Chapter 13	Contents
<ul><li>Proc</li><li>Serv</li></ul>	parison ess
13.2 Digital S	ignature Schemes
13.3 Summar	У
	13.4



hapter 13		13.1 Introduction
		Comparis
	Conventional Signature	Digital Signatures
(1) Inclusion	Included in the document as part of the document.	Send the signature as a separate document.
(2) Verification Method	Recipient compares the signature on the document with the signature on file.	<ul> <li>The recipient receives the message and the signature.</li> <li>The recipient needs to apply a verification technique to the combination of the message and the signature to verify the authenticity.</li> </ul>
(3) Relationship	Normally a one-to-many relationship between a signature and documents.	<i>One-to-one</i> relationship between a signature and a message.
(4) Duplicity	A copy of the signed document can be distinguished from the original one on file.	No such distinction unless there is a factor of time on the document



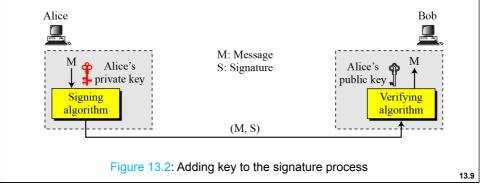


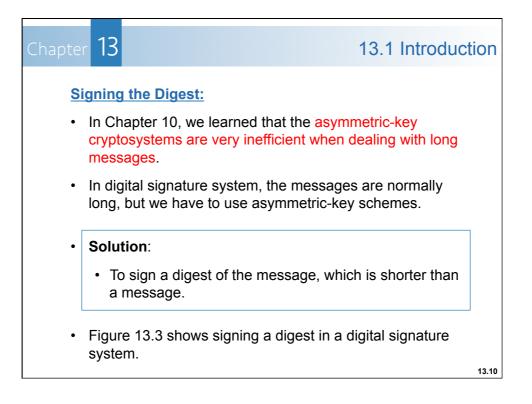


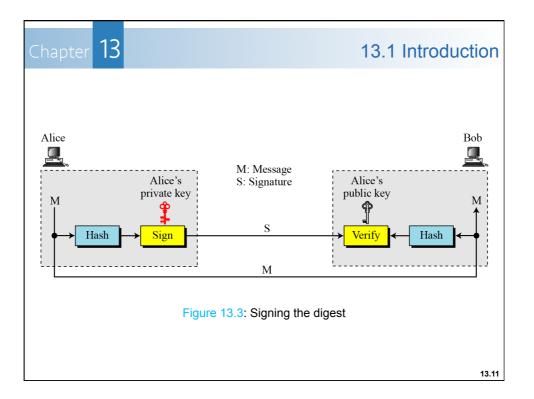
## Need for Keys:

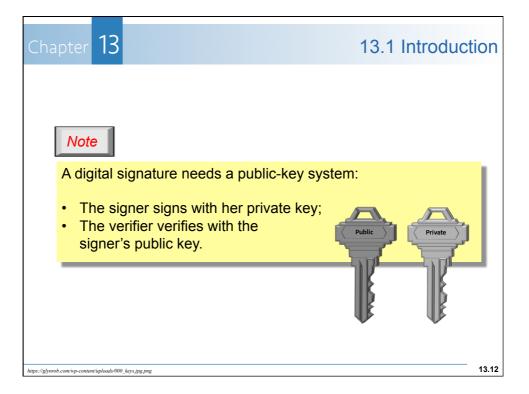
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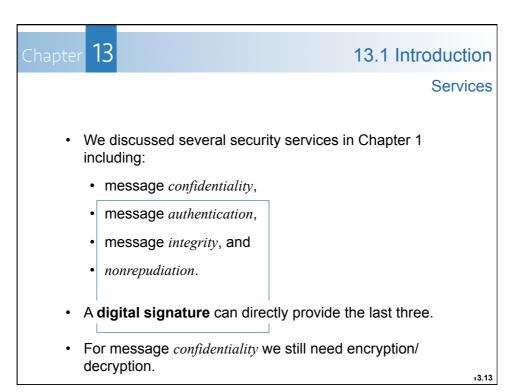
- In digital signature, the signer uses his private key, applied to a *signing algorithm*, to **sign** the document.
- The verifier (recipient), uses the public key of the signer, applied to the *verifying algorithm*, to **verify** the document.

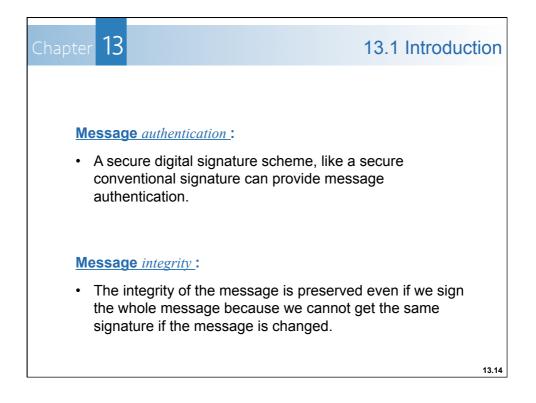


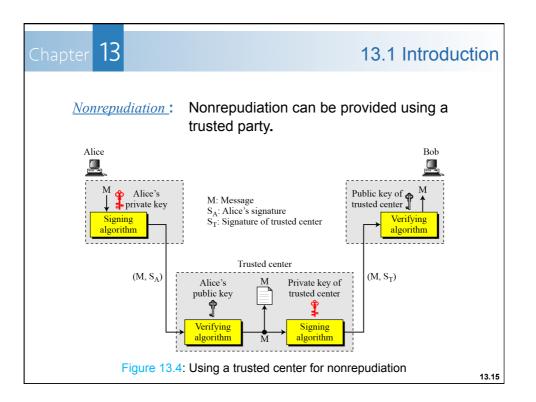


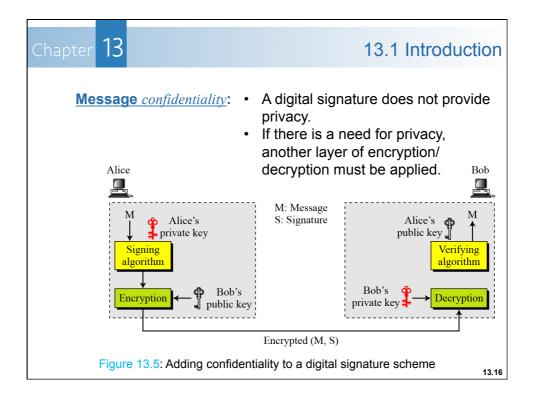


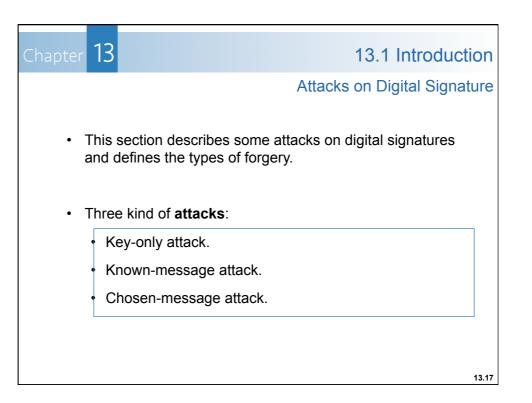


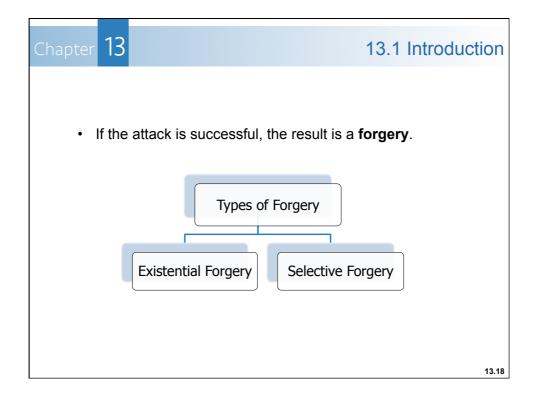


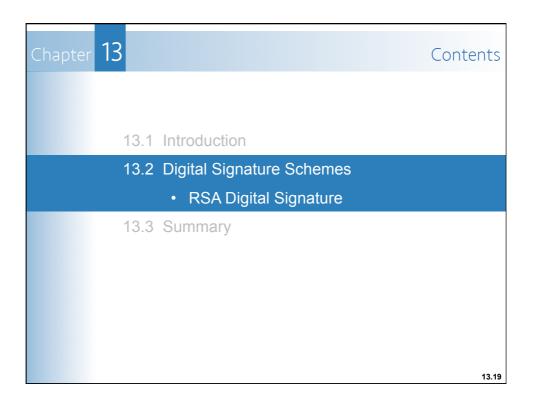


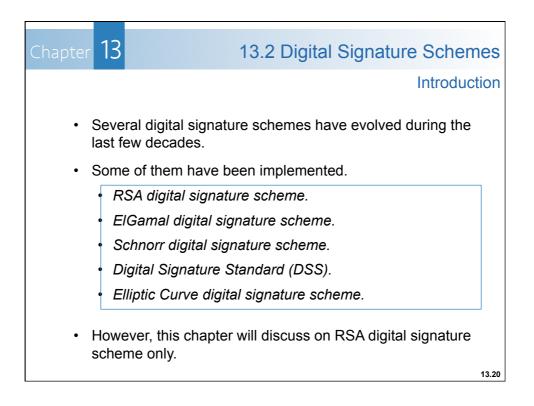




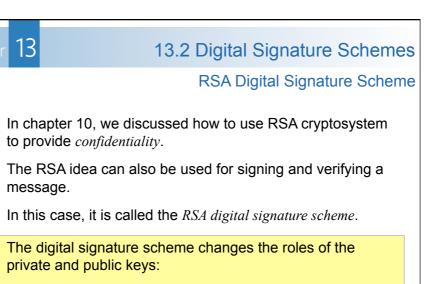




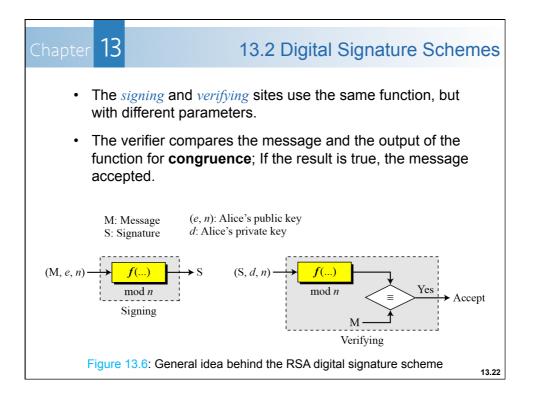




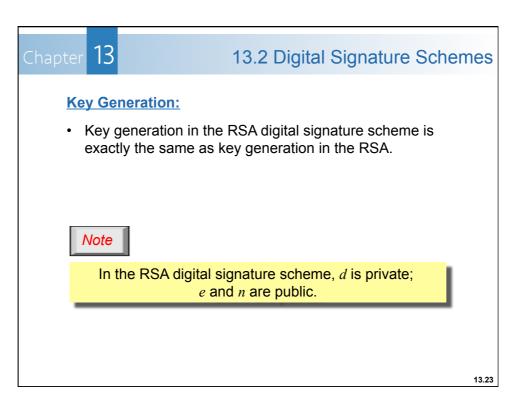
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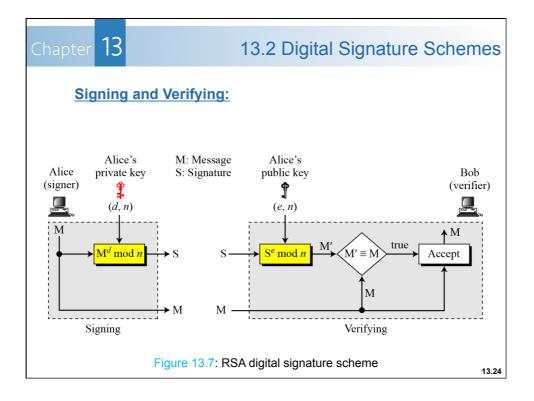


- ① The private and public keys of the senders are used.
- ② The sender uses his own private key to sign the document; the receiver uses the sender's public key to verify the document.

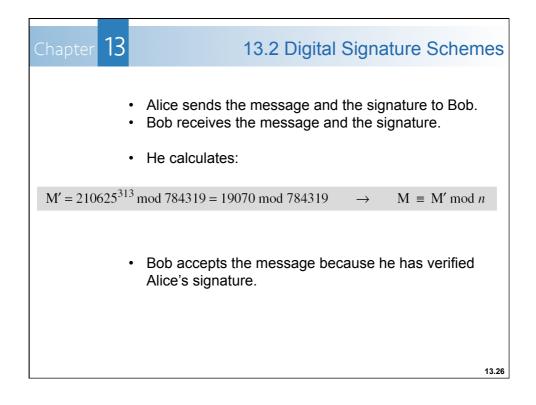


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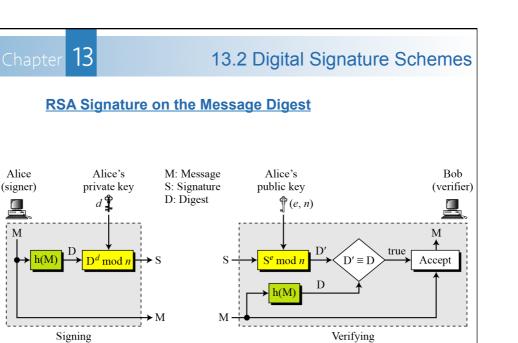




Chapter 13	13.2 Digital Signature Schemes
	s a trivial example, suppose that Alice chooses $p = 823$ and $q = 953$ , and calculates $n = 784319$ .
•	The value of $\phi(n)$ is 782544. Now she chooses $e = 313$ and calculates $d = 160009$ .
	At this point key generation is complete. Now imagine that Alice wants to send a message with the value of M = 19070 to Bob. She uses her private exponent, e = 160009, to sign the message:
M: 19070	$\rightarrow$ S = (19070 <sup>160009</sup> ) mod 784319 = 210625 mod 784319
	(continued)



Signing



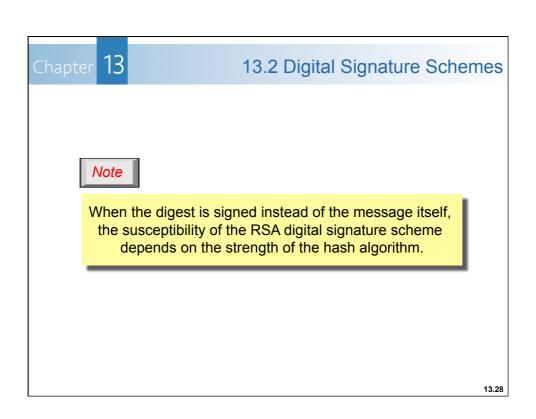
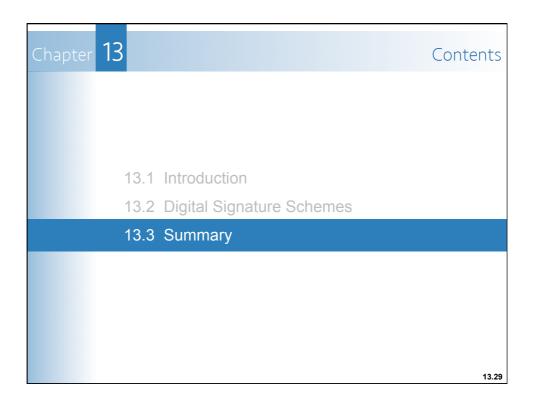
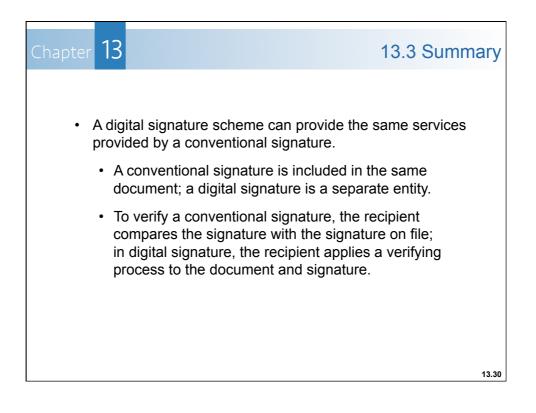
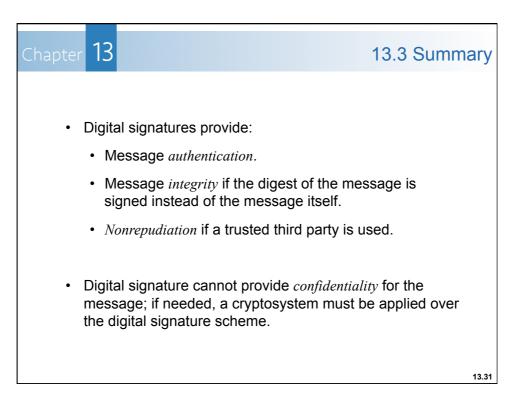


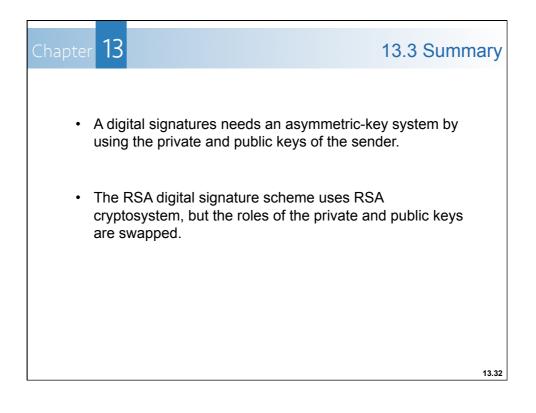
Figure 13.8: The RSA signature on the message digest

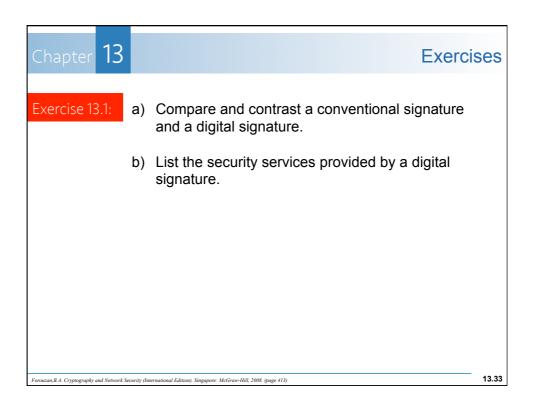
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Chapter 13	Exercis	ses
	<ul> <li>Using the RSA scheme, let p = 809, q = 751, and d = 23 Calculate the public key <i>e</i>. Then,</li> <li>a) Sign and verify a message with M<sub>1</sub> = 100. Call the signature S<sub>1</sub>.</li> <li>b) Sign and verify a message with M<sub>2</sub> = 50. Call the signature S<sub>2</sub>.</li> </ul>	3.
Forouzan, B.A. Cryptography and Network Se	curity (International Edition). Singapore: McGraw-Hill, 2008. (page 413)	13.34